Algorithms

Shortest Paths

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- Let G be a weighted graph. The **length** (or weight) of a path, P, is the sum of the weights of the edges of P.
- The distance from a vertex v to a vertex u in G, denoted d(u,v) is the length of a minimum length path(also called shortest path) from v to u, if such a path exists.
- Given as input a weighted graph, G = (V,E), and a distinguished vertex, s, we want to find the shortest weighted path from s to every other vertex in G. This is known as the single source shortest paths problem.
- One algorithm which can be used to compute the shortest path is Dijkstra's algorithm.

- Dijkstra's algorithm is another example of a greedy algorithm
- It generates the shortest paths in stages.
- In each stage, a shortest path to a new destination vertex is generated
- The destination for the next shortest path is selected using the greedy criterion: From the vertices to which a shortest path has not been generated, select one that results in a least path length.

- We begin with the trivial path from the source vertex to itself. This path has no edges and has a length of 0.
- In each stage of the greedy algorithm, the next shortest path is generated. This next shortest path is the shortest possible one edge extension of an already generated shortest path.



- Using the above graph we choose v_1 as our starting vertex.
- We now make this node the reference node and mark it as known.
- We maintain a table of distances d_v and vertices which cause a change to d_v which we call p_v

Initial table configuration

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	No	0	0
V ₂	No	-	0
V ₃	No	-	0
V_4	No	-	0
V_5	No	-	0
V_6	No	-	0
V ₇	No	-	0

 We now calculate the distance from v₁ to all of its adjacent nodes which are not known and update the table

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V ₂	No	2	V ₁
V ₃	No	4	V ₁
V_4	No	1	V ₁
V_5	No	-	0
V ₆	No	-	0
V ₇	No	-	0

• We now mark v₄ as known and recalculate the table

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V_2	No	2	V ₁
V_3	No	3	V_4
V_4	Yes	1	V ₁
V_5	No	3	V_4
V_6	No	9	V_4
V ₇	No	5	V_4

• We now mark v₂ as known and recalculate the table

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V ₂	Yes	2	V ₁
V ₃	No	3	V_4
V_4	Yes	1	V ₁
V_5	No	3	V_4
V ₆	No	9	V_4
V ₇	No	5	V_4

• We now mark v₃ as known and recalculate the table

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V ₂	Yes	2	V ₁
V ₃	Yes	3	V_4
V ₄	Yes	1	V ₁
V_5	No	3	V_4
V ₆	No	8	V_3
V ₇	No	5	V_4

• We now mark v₅ as known and recalculate the table

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V ₂	Yes	2	V ₁
V ₃	Yes	3	V_4
V_4	Yes	1	V ₁
V_5	Yes	3	V_4
V ₆	No	8	V_3
V ₇	No	5	V_4

• We now mark v₇ as known and recalculate the table

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V ₂	Yes	2	V ₁
V ₃	Yes	3	V_4
V_4	Yes	1	V ₁
V_5	Yes	3	V_4
V ₆	No	6	V_7
V ₇	Yes	5	V_4

• We now mark v₆ as known

vertex	known	d_{v}	\boldsymbol{p}_{v}
V ₁	Yes	0	0
V ₂	Yes	2	V ₁
V_3	Yes	3	V_4
V_4	Yes	1	V ₁
V_5	Yes	3	V_4
V ₆	Yes	6	V_7
V ₇	Yes	5	V_4











